EXPANDING GRAPHS CONTAIN ALL SMALL TREES

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Received 2 May 1986

The assertion of the title is formulated and proved. The result is then used to construct graphs with a linear number of edges that, even after the deletion of almost all of their edges or almost all of their vertices, continue to contain all small trees.

1. Introduction

If H is an undirected graph, V(H) will denote its set of vertices and E(H) will denote its set of edges. If $X \subseteq V(H)$, $\Gamma_H(X)$ will denote the set of neighbors in H of vertices in X. If X is a set, |X| will denote its cardinality.

The following theorem, which is implicit in a result of Pósa [5], has been given an elegant proof by Lovász ([3], Ch. 10. Problem 20).

Theorem 0. If H is a non-empty graph such that, for each $X \subseteq V(H)$ with $|X| \le n$,

$$|\Gamma_H(X) \setminus X| \ge 2|X|-1,$$

then H contains a path with 3n-2 vertices.

Using Theorem 0, Beck [2] proved an upper bound of the form O(n) for the minimum possible number of edges in graphs that, even after the deletion of half their edges, continue to contain a path with n vertices; Alon and Chung [1] have given an explicit construction for such graphs.

The main result of this paper is the following theorem.

Theorem 1. If H is a non-empty graph such that, for every $X \subseteq V(H)$ with $|X| \le 2n-2$,

$$|\Gamma_H(X)| \ge (d+1)|X|,$$

then H contains every tree with n vertices and maximum degree at most d.

Since a path with *n* vertices is the unique tree with *n* vertices and maximum degree 2, Theorem 1 generalizes the essence of Theorem 0 from paths to trees. Using Theorem 1, the arguments of Beck [2] and Alon and Chung [1], and a recent result of Lubotzky, Phillips and Sarnak [4], we prove the following theorem.

Theorem 2. Let $\delta > 0$ and d be fixed. For every n there is a graph F with O(n) edges that, even after deletion of all but $\delta |E(F)|$ edges, continues to contain every tree with n vertices and maximum degree at most d.

For $\delta = 1/2$, Beck [2] proved an upper bound, without an explicit construction, of the form $O(n(\log n)^{12})$. We shall also prove the following theorem.

Theorem 3. Let $\varepsilon > 0$ and d be fixed. For every n there is a graph F with O(n) edges that, even after deletion of all but $\varepsilon |V(F)|$ vertices, continues to contain every tree with n vertices and maximum degree at most d.

2. Proof of Theorem 1

If T is a tree and H is a graph, a map $f: V(T) \rightarrow V(H)$ will be called an *embedding* of T in H if it is injective and f(v) and f(w) are adjacent in H whenever v and w are adjacent in T. A tree T will be called *small* if it has at most n vertices and maximum degree at most d. (The parameters n and d will remain fixed throughout this proof.) A graph H will be called *expanding* if, for every $X \subseteq V(H)$ with $|X| \le 2n-2$,

$$|\Gamma_H(X)| \cong (d+1)|X|.$$

Our goal is to show that if T is small and H is non-empty and expanding, then there is an embedding of T in H. To achieve this we shall define a class of "good" embeddings. We shall then show that this class has the following two properties.

Property 1. If T consists of a single vertex and H is a non-empty expanding graph, then there is a good embedding of T in H.

Property 2. If T is a small tree and S is a subtree of T obtained by deleting a leaf and the edge incident with it, then any good embedding of S in an expanding graph H can be extended to a good embedding of T in H.

When this has been done, it will follow by induction on |T| that, if T is a small tree and H is a non-empty expanding graph, then there is a good embedding of T in H. If |V(T)|=1, this follows from Property 1. If $|V(T)| \ge 2$, let S be any tree obtained from T by deleting a leaf and the edge incident with it. By inductive hypothesis, there is a good embedding of S in H, and by Property 2, this can be extended to a good embedding of T in H. This completes the induction and the proof of Theorem 1.

To define good embeddings, we shall need some auxiliary definitions. Let f be an embedding of a tree T in a graph H. If $X \subseteq V(H)$, we shall define the assets $A_f(X)$ of X under f to be $|\Gamma_H(X) \setminus f(V(T))|$. If $x \in V(H)$, we shall let $J_f(x)$ denote the degree of $f^{-1}(x)$ in T if $x \in f(T)$, and 0 otherwise. We shall let $B_f(x)$ denote $d-J_f(x)$. If $X \subseteq V(H)$, we shall define the liabilities $B_f(X)$ of X under f to be $\sum_{x \in X} B_f(x)$, and the balance $C_f(X)$ of X under f to be $A_f(X) - B_f(X)$. A set $X \subseteq V(H)$ will be called solvent under f if $C_f(X) \ge 0$, critical under f if $C_f(X) = 0$, and bankrupt under f if $C_f(X) < 0$. Finally, we arrive at the key definition. An embedding f of a tree T in a graph H is good if every $X \subseteq V(H)$ with $|X| \le 2n-2$ is solvent. It remains to prove Properties 1 and 2.

To prove Property 1, suppose that T is a tree consisting of a single vertex and that H is a non-empty expanding graph. Since H is non-empty, there is an embedding f of T in H. We shall show that f is good. Suppose $X \subseteq V(H)$ and $|X| \le 2n-2$. Since H is expanding,

$$|\Gamma_H(X)| \ge (d+1)|X|.$$

Since f(V(T)) consists of a single vertex,

$$A_f(X) = |\Gamma_H(X) \setminus f(V(T))| \ge d|X|.$$

On the other hand, $J_f(x) = 0$, so $B_f(x) = d$ and

$$B_t(X) = d|X|.$$

Thus X is solvent, which completes the proof of Property 1.

To prove Property 2, suppose that T is a small tree, that S is a subtree obtained from T by deleting a leaf v and the incident edge $\{v, w\}$, and that f is a good embedding of S in an expanding graph H. Let G denote the set of embeddings of T in H that are extensions of f. Let $Y = \Gamma_H(\{f(w)\}) \setminus f(V(S))$. The map $g \mapsto g(v)$ is clearly a bijection between G and Y. We must show that some $g \in G$ is good.

Suppose, to obtain a contradiction, that no $g \in G$ is good. Then for every $g \in G$, there is a set $X_g \subseteq V(H)$ with $|X_g| \cong 2n-2$ and X_g bankrupt under g. On the other hand, since f is a good embedding, we have X_g solvent under f for all $g \in G$. Since $g(V(T)) = f(V(S)) \cup \{g(v)\}$, we have

$$A_g(X) = A_f(X) - I(g(v), \Gamma_H(X)),$$

where I(z, Z) is 1 if $z \in Z$ and 0 otherwise. Since $E(T) = E(S) \cup \{\{v, w\}\}\$, we have

$$B_g(X) = B_f(X) - I(f(w), X) - I(g(v), X).$$

Thus,

$$C_g(X) = C_f(X) - I(g(v), \Gamma_H(X)) + I(f(w), X) + I(g(v), X).$$

For each $g \in G$ we have $C_g(X_g) < 0$ and $C_f(X_g) \ge 0$, so we must have $C_f(X_g) = 0$, $I(g(v), \Gamma_H(X)) = 1$, I(f(w), X) = 0 and I(g(v), X) = 0. Thus for every $g \in G$, we have X_g critical, $g(v) \in \Gamma_H(X_g)$, $f(w) \notin X_g$ and $g(v) \notin X_g$. To proceed further, we shall need some lemmas.

Lemma 1.1. If $X \subseteq V(H)$ is critical under f and $|X| \le 2n-2$, then $|X| \le n-1$.

Proof. Since H is an expanding graph and $|X| \le 2n-2$,

$$|\Gamma_H(X)| \ge (d+1)|X|.$$

Since T is a small tree and S is obtained by deleting a leaf, $|V(S)| \le n-1$, and thus

$$A_f(X) = |\Gamma_H(X) \setminus f(V(S))| \ge (d+1)|X| - (n-1).$$

On the other hand, $J_f(x) \ge 0$, so $B_f(x) \le d$ and

$$B_{\varepsilon}(X) \leq d|X|.$$

Since X is critical under f, $A_f(X) = B_f(X)$, so $|X| \le n-1$.

Lemma 1.2. The balance $C_f(\cdot)$ is submodular; that is

$$C_f(X \cup Y) + C_f(X \cap Y) \leq C_f(X) + C_f(Y).$$

Proof. Since $B_f(X)$ is defined by a sum over vertices in X, $B_f(\cdot)$ is modular; that is

$$B_f(X \cup Y) + B_f(X \cap Y) = B_f(X) + B_f(Y).$$

Thus it suffices to show that $A_f(\cdot)$ is submodular; but this is an immediate consequence of the relations $\Gamma_H(X \cup Y) = \Gamma_H(X) \cup \Gamma_H(Y)$ and $\Gamma_H(X \cap Y) \subseteq \Gamma_H(X) \cap \Gamma_H(Y)$.

Lemma 1.3. If $X, Y \subseteq V(H)$ are critical under f and $|X|, |Y| \le n-1$, then $X \cup Y$ is critical under f and $|X \cup Y| \le n-1$.

Proof. Since f is a good embedding and $|X \cup Y|$, $|X \cap Y| \le 2n-2$, $C_f(X \cup Y)$, $C_f(X \cap Y) \ge 0$. Since X and Y are critical under f, $C_f(X)$, $C_f(Y) = 0$. Thus Lemma 1.2 implies $C_f(X \cup Y) \le 0$, whence $C_f(X \cup Y) = 0$ and $X \cup Y$ is critical. Lemma 1.1 now implies $|X \cup Y| \le n-1$.

We now resume the proof of Property 2. For every $g \in G$, X_g is critical under f and $|X_g| \le 2n-2$; thus by Lemma 1.1, $|X_g| \le n-1$. We now claim that the set

$$X^* = \bigcup_{g \in G} X_g$$

is critical under f and that $|X^*| \le n-1$. If |G| = 0, this is trivial; if $|G| \ge 1$, it follows by induction using Lemma 1.3. Now consider $X' = X^* \cup \{f(w)\}$. Since f is good and $|X'| \le n$, X' is solvent under f. Since $g(v) \in \Gamma_H(X_g)$ for every $g \in G$, $Y \subseteq \Gamma_H(X^*)$, which implies

$$A_f(X') = A_f(X^*).$$

Since $f(w) \notin X_g$ for every $g \in G$, $f(w) \notin X^*$, which implies

$$B_f(X') = B_f(X^*) + B_f(f(w)).$$

Since T is small, w has degree at most d in T and thus degree at most d-1 in S, so $J_f(f(w)) \le d-1$ and $B_f(f(w)) \ge 1$. Thus $B_f(X') > B_f(X^*)$, so $C_f(X') < C_f(X^*)$. Since X^* is critical under f, X' is bankrupt under f. This contradicts the fact that X' is solvent under f, completing the proof of Property 2 and of Theorem 1.

3. Proof of Theorem 2

Let us say that a graph F with m vertices is (k, l)-spectral if

(A) every vertex of F has degree k, so that the adjacency matrix of F (which, by abuse of notation, we shall also denote by F), has an eigenvalue $\lambda_0(F) = k$ corresponding to an eigenvector e that assigns the value 1 to every vertex of F; and

(B) each of the other eigenvalues $\lambda_j(F)$, $1 \le j \le m-1$, of F satisfies $|\lambda_j(F)| \le l$. Lubotzky, Phillips and Sarnak [4] have shown that if p and q are primes congruent to 1 modulo 4, with p a quadratic non-residue modulo q, then there is an explicitly construced graph F that

- (1) has $m = q(q^2 1)/2$ vertices;
- (2) is $(k, 2(k-1)^{1/2})$ -spectral, where k=p+1; and
- (3) has girth at least $(2/3) \log_{k-1} m$.

If F is a graph and $X \subseteq V(F)$, $\Theta_F(X)$ will denote the set of edges of F having both ends in X, $\Phi_F(X)$ will denote the set of edges of F having at least one end in X, and $\Psi_F(X)$ will denote the set of edges of F having exactly one end in X. We shall need the following lemma.

Lemma 2.1. If a graph F has m vertices and is (k, l)-spectral, and if $X \subseteq V(F)$, then

$$|\Theta_F(X)| \leq k|X|^2/2m + l|X|/2.$$

Proof. Define the function g on V(F) by g(x)=1-|X|/m if $x \in X$ and g(x)==-|X|/m otherwise. Then

$$|\Psi_F(X)| = \frac{1}{2} \sum_{\{x,y\} \in E(F)} (g(x) - g(y))^2.$$

Expanding the sum yields

$$\begin{aligned} |\Psi_F(X)| &= k \sum_{\mathbf{x} \in V(F)} g(\mathbf{x})^2 - \sum_{\{\mathbf{x}, \mathbf{y}\} \in E(F)} g(\mathbf{x}) g(\mathbf{y}) \\ &= k \langle g, g \rangle - \langle g, Fg \rangle, \end{aligned}$$

where $\langle \cdot, \cdot \rangle$ denotes inner product. Since g is orthogonal to the eigenvector e, $\langle g, Fg \rangle \leq l \langle g, g \rangle$. Thus $|\Psi_F(X)| \geq (k-l) \langle g, g \rangle = (k-l) |X| (1-|X|/m)$. Since $|\Theta_F(X)| = (k|X|-|\Psi_F(X)|)/2$, $|\Theta_F(X)| \leq k|X|^2/2m+l|X|(1-|X|/m)/2 \leq k|X|^2/2m+l|X|/2$.

Let $\delta > 0$, d and n be given. Assign p and q the smallest appropriate values for which the following conditions are satisfied.

Condition 1. $k \ge 16(d+2)^2/\delta^2$.

Condition 2. $m \ge 2(d+2)^2(2n-2)/\delta$.

If δ and d are fixed, then k=O(1) and, by the prime number theorem for arithmetic progressions, m=O(n). Thus F has km/2=O(n) edges. To complete the proof of Theorem 2, we shall show that, if $D\subseteq E(F)$ with $|D| \ge \delta km/2$, then the graph G with vertices V(G)=V(F) and edges E(G)=D contains an expanding graph.

Let $W \subseteq V(G)$ be a minimal non-empty set such that $|\Theta_G(W)| \ge \delta k |W|/2$ (such a set exists, since $|\Theta_G(V(G))| \ge \delta k |V(G)|/2$). Let H be the graph with vertices V(H) = W and edges $E(H) = \Theta_G(W)$. We shall show that H is an expanding graph.

First, we claim that if $X \subseteq V(H)$, then $|\Phi_H(X)| \ge \delta k|X|/2$. This is clear if $X = \emptyset$ or X = W. If $\emptyset \subset X \subset W$, then $W \setminus X$ is a proper non-empty subset of W, so by the minimality of W, $|\Theta_H(W \setminus X)| = |\Theta_G(W \setminus X)| < \delta k|W \setminus X|/2$. Thus $|\Phi_H(X)| = |\Theta_H(W)| - |\Theta_H(W \setminus X)| > \delta k|W|/2 - \delta k|W \setminus X|/2 = \delta k|X|/2$, and the claim is proved.

Now suppose that $X \subseteq V(H)$ with $|X| \le 2n-2$. We shall show that $|\Gamma_H(X)| \ge (d+1)|X|$. Suppose, to obtain a contradiction, that $|\Gamma_H(X)| < (d+1)|X|$. Let $Y = X \cup \Gamma_H(X)$. Then $|Y| < (d+2)|X| \le (d+2)(2n-2)$. By Lemma 2.1 (with l = 1)

 $=2(k-1)^{1/2} \le 2k^{1/2})$ we have $|\Theta_H(Y)| \le |\Theta_F(Y)| \le k|Y|^2/2m + k^{1/2}|Y|$. By supposition we have $|\Theta_H(Y)| \le k(d+2)^2(2n-2)|X|/2m + k^{1/2}|Y|$, and by Condition 2, $|\Theta_H(Y)| \le \delta k|X|/4 + k^{1/2}|Y|$. Since $\Phi_H(X) \subseteq \Theta_H(Y)$, by the claim we have $|\Theta_H(Y)| \ge \delta k|X|/2$. Thus $k^{1/2}|Y| \ge \delta k|X|/4$, so $|Y| \ge \delta k^{1/2}|X|/4$. By Condition 1, $k^{1/2} \ge 4(d+2)/\delta$, so $|Y| \ge (d+2)|X|$. This contradiction shows that H is an expanding graph and completes the proof of Theorem 2.

4. Proof of Theorem 3

We shall need the following lemma.

Lemma 3.1. If a simple graph F has m vertices and is (k, l)-spectral, and if $X \subseteq V(F)$, then

$$|\Theta_F(X)| \ge k|X|^2/2m - (l+2)|X|/2.$$

Proof. Since F is a simple graph, so is its complement \overline{F} , which has adjacency matrix J-I-F, where J is an $m\times m$ matrix with every entry equal to 1 and I is an $m\times m$ identity matrix. The graph \overline{F} has eigenvalues $\lambda_0(\overline{F})=m-1-k$ and $\lambda_j(\overline{F})=-1-\lambda_j(F)$, $1\leq j\leq m-1$. Thus \overline{F} is (m-1-k,l+1)-spectral. Applying Lemma 2.1 to $|\Theta_F(X)|$ in the formula $|\Theta_F(X)|=|X|(|X|-1)/2-|\Theta_F(X)|$ completes the proof.

Let $\varepsilon > 0$, d and n be given. Set $\delta = \varepsilon^2/2$. Assign p and q the smallest appropriate values for which Conditions 1 and 2 in the proof of Theorem 2 are satisfied and for which $m \ge k^3$.

If ε and d are fixed, so is δ . Thus again k=O(1) and m=O(n), and the graph F has km/2=O(n) edges. Furthermore, $m \ge k^3$ ensures that the girth of F exceeds 2, so F is a simple graph. To complete the proof of Theorem 3, we shall show that, if $U \subseteq V(F)$ with $|U| \ge \varepsilon m$, then $|\Theta_F(U)| \ge \delta km/2$, for then we may set $D = \Theta_F(U)$ and apply Theorem 2.

Suppose, to obtain a contradiction, that $U \subseteq V(F)$ with $|U| \cong \epsilon m$ but $|\Theta_F(U)| < \delta km/2$. Since $\delta = \epsilon^2/2$, $|\Theta_F(U)| < \epsilon^2 km/4 \le k|U|^2/4m$. By Lemma 3.1 (with $l+2=2(k-1)^{1/2}+2\le 4k^{1/2}$) we have $|\Theta_F(U)| \ge k|U|^2/2m-2k^{1/2}|U|$. Thus $2k^{1/2}|U| > k|U|^2/4m$, so $|U| < 8m/k^{1/2}$. By Condition 1, $k^{1/2} \ge 4(d+2)/\delta \ge 8/\epsilon$, so $|U| < \epsilon m$. This contradiction completes the proof of Theorem 3.

5. References

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